Simulating and Verifying Cyber-Physical Systems: Current Challenges and Novel Research Directions

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Discrete Models

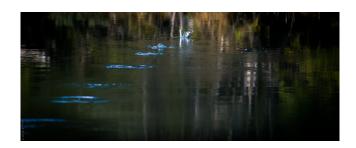
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Continuous Models



Hybrid Models



Hybrid Models: Discrete ∪ Continuous

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Computer Science (Automata Theory)

Essentially **discrete**: finite set of modes with continuous evolution within modes. [Hybrid Automata, Alur et al. 1992]

Control (Differential Equations)

Essentially **continuous**: non-smooth (discontinuous) dynamics, differential inclusions, Filippov/Utkin regularization. [Mosterman 1998, Sanfelice 2003]

Define/Describe Reactions to Events

Events

Model a simplified (often discrete) perception of the rich environment.

Reactions

Model how the system is supposed to react to an event so that it respects a set of constraints which are essentially physics laws and/or predefined requirements.

Convenient Model for a Large Class of Systems









Modeling is Instrumental to Master Complexity

Abstraction/Refinement Relationships

The molecular composition of the stratosphere has a minor impact on the Earth's orbit. Likewise, a galaxy is a dot in the Laniakea supercluster.

Compositionality and Reuse

Human beings are perhaps the most extreme example of both concepts: nature builds on top of what works to create new more complex structures.

Simulation (Time Travel)

A peek in the future (or the past) of a concrete model for a concrete initial (or final) condition. Essentially by "executing" the model step by step: only a **Local** recipe is needed.

Verification (Time Abstraction)

Qualitative analysis of the geometry (shape) of the state space as it captures **Global** properties.

Local description defines the global properties which in turn benefit numerical approximations (eg. Geometrical Integration)

Outline

- 1 Introduction
- 2 Verification
- 3 Simulation
- 4 Challenges

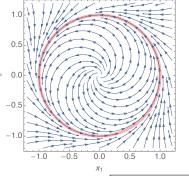
Global Properties

If one regards a sorting algorithm as a discrete dynamical system acting on the given set, then the sorted list is an **invariant** or fixed point. It is an attractor that is reached from any initial position in finite time (# steps).



Qualitative Analysis of Dynamical Systems

$$(\dot{x_1}, \dot{x_2}) = (x_1 - x_1^3 - x_2 - x_1x_2^2, x_1 + x_2 - x_1^2x_2 - x_2^3)$$



Invariant Equation

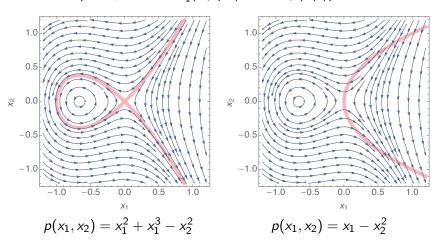
Algebraic

The solution for $\mathbf{x}_0 = (1,0)$ respects $|x_1(t)|^2 + x_2(t)^2 - 1 = 0$

$$x_1(t)^2 + x_2(t)^2 - 1 = 0$$

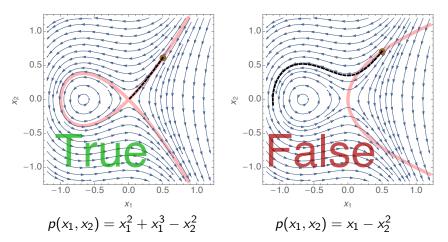
Problem I. Checking Invariance of Algebraic Equations

Given
$$\dot{\mathbf{x}} = (-2x_2, -2x_1 - 3x_1^2), \ p(\mathbf{x}_0) = 0$$
, is $p(\mathbf{x}(t)) = 0$ for all t ?



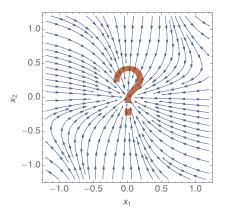
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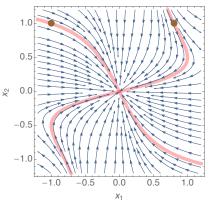
Problem II. Generate Algebraic Invariant Equations

Given $\dot{\mathbf{x}} = (-x_1 + 2x_1^2x_2, -x_2)$, how to generate p such that $p(\mathbf{x}(t)) = 0$?



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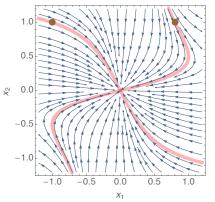
Given $\dot{\mathbf{x}} = (-x_1 + 2x_1^2x_2, -x_2)$, how to generate p such that $p(\mathbf{x}(t)) = 0$?



$$p_{(x_1(0),x_2(0))}(x_1,x_2) = (x_2(0) - x_1(0)x_2(0)^2)x_1 - x_1(0)(x_2 - x_1x_2^2) = 0$$

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$$\begin{split} p_{(x_1(0),x_2(0))}\big(x_1,x_2\big) &= (x_2(0)-x_1(0)x_2(0)^2)x_1 - x_1(0)\big(x_2-x_1x_2^2\big) = 0 \\ &\frac{x_1}{x_2-x_1x_2^2} \text{ is an invariant } \textbf{rational function}. \end{split}$$

Gradient
$$\nabla p := (\frac{\partial p}{\partial x_1}, \dots, \frac{\partial p}{\partial x_n})$$

Lie Derivation
$$\mathfrak{D}_{\mathbf{f}}(p) := \frac{dp(\mathbf{x}(t))}{dt} = \langle \nabla p, \mathbf{f} \rangle$$
 $(\dot{\mathbf{x}} = \mathbf{f})$

Singular Locus

$$\mathsf{SL}(p) := \left\{ \mathbf{x} \in \mathbb{R}^n \mid \nabla p = \mathbf{0} \right\} = \left\{ \mathbf{x} \in \mathbb{R}^n \mid \frac{\partial p}{\partial x_1} = 0 \land \dots \land \frac{\partial p}{\partial x_n} = 0 \right\}$$

 $\mathbf{x} \in V_{\mathbb{R}}(p)$ $(p(\mathbf{x}) = 0)$ is singular if $\mathbf{x} \in SL(p)$, regular otherwise.

Notation for " $S \subseteq \mathbb{R}^n$ is invariant for \mathbf{f} "

$$S \rightarrow [\dot{\mathbf{x}} = \mathbf{f}]S$$

The set S is an invariant set for \mathbf{f}

Starting with \mathbf{x}_0 s.t $\mathbf{x}_0 \in S$: for all t > 0, $\mathbf{x}(t)$ solution of the IVP $(\dot{\mathbf{x}} = \mathbf{f}, \mathbf{x}(0) = \mathbf{x}_0)$ is in S

N.B. Treating $\dot{\mathbf{x}} = \mathbf{f}$ as a program, one can think of the top formula as representing the Hoare triple $\{S\}$ $\dot{\mathbf{x}} = \mathbf{f}$ $\{S\}$.

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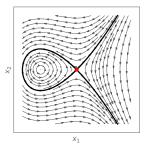
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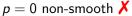
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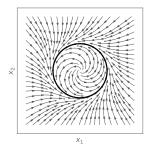
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Necessary and sufficient for smooth invariant manifolds (Lie, 1893).

$$\text{(Lie)} \frac{\rho = 0 \to (\mathfrak{D}_{\mathbf{f}}(\rho) = 0 \land \nabla \rho \neq \mathbf{0})}{(\rho = 0) \to [\dot{\mathbf{x}} = \mathbf{f}] \ (\rho = 0)}$$







$$p=0$$
 smooth \checkmark

No flow in the problem variables at singularities on the variety

$$(\mathsf{Lie}^{\circ})\frac{p=0\to \big(\mathfrak{D}_{\mathbf{f}}(p)=0\land \big(\nabla p=\mathbf{0}\to \mathbf{f}=\mathbf{0}\big)\big)}{(p=0)\to [\dot{\mathbf{x}}=\mathbf{f}]\ (p=0)}$$

Flow at singularities on the variety is directed into the variety

$$(\mathsf{Lie}^*) \frac{p = 0 \to (\mathfrak{D}_{\mathbf{f}}(p) = 0 \land (\nabla p = \mathbf{0} \to p(\mathbf{x} + \lambda \mathbf{f}) = 0))}{(p = 0) \to [\dot{\mathbf{x}} = \mathbf{f}] \ (p = 0)}$$

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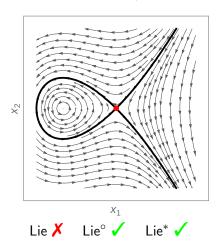
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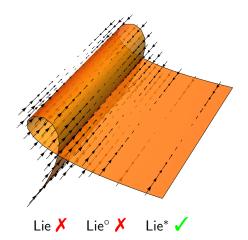
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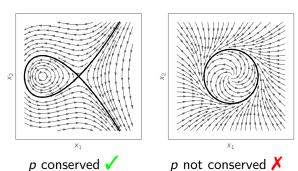




[Platzer, J. Log. Comput. 2010]

Necessary and sufficient for conserved quantities (integrals of motion).

$$(\mathsf{FI})\frac{\mathfrak{D}_{\mathbf{f}}(p) = 0}{(p = 0) \to [\dot{\mathbf{x}} = \mathbf{f}] \ (p = 0)}$$



Continuous consecutions (C-c) and **polynomial consecutions** (P-c) are Darboux polynomials (Darboux, 1878).

$$(\mathsf{C-c})\frac{\exists \lambda \in \mathbb{R}, \ \mathfrak{D}_{\mathbf{f}}(p) = \lambda p}{(p=0) \to [\dot{\mathbf{x}} = \mathbf{f}] \ (p=0)},$$

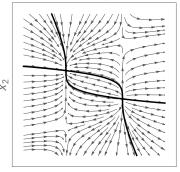
$$(P-c)\frac{\exists \lambda \in \mathbb{R}[\mathbf{x}], \ \mathfrak{D}_{\mathbf{f}}(p) = \lambda p}{(p=0) \to [\dot{\mathbf{x}} = \mathbf{f}] \ (p=0)} \ .$$

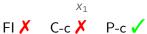
Extensions of FI

[Sankaranarayanan et al., FMSD 2008]

$$\mathbf{f} = (3(x_1^2 - 4), -x_2^2 + x_1x_2 + 3), \qquad p = x_2^4 + 2x_1x_2^3 + 6x_2^2 + 2x_1x_2 + x_1^2 + 3,$$

$$\mathfrak{D}_{\mathsf{f}}(p) = \underbrace{\left(6x_1 - 4x_2\right)}_{\lambda} p$$





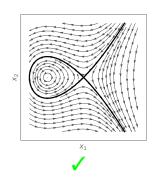


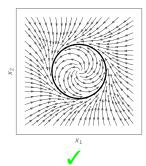
Differential Radical Invariants (DRI)

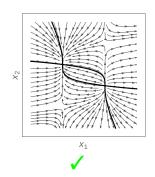
[G. et al., TACAS 2014, SAS 2014]

Necessary and sufficient for invariant varieties.

$$(\mathsf{DRI})\frac{\rho=0\to \bigwedge_{i=0}^{N-1}\mathfrak{D}_{\mathbf{f}}^{(i)}(\rho)=0}{(\rho=0)\to [\dot{\mathbf{x}}=\mathbf{f}]\ (\rho=0)}$$







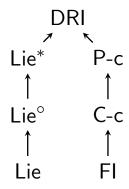
$$(\mathsf{R}_A) \frac{A}{S_A : T_A \longrightarrow [\dot{\mathbf{x}} = \mathbf{f}] S_A : T_A} \qquad (\mathsf{R}_B) \frac{B}{S_B : T_B \longrightarrow [\dot{\mathbf{x}} = \mathbf{f}] S_B : T_B}$$

Partial Order

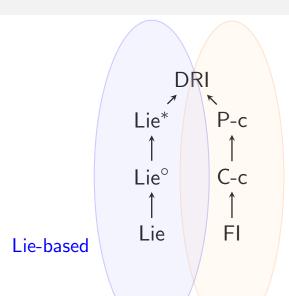
 $R_A \preccurlyeq R_B$ if and only if $A \implies B$ and T_A is a "subtype" of T_B .

- $R_A \sim R_B \ (R_A \leq R_B \text{ and } R_A \geq R_B)$ Equivalence.
- $R_A \prec R_B \ (R_A \preccurlyeq R_B \ \text{and} \ R_A \not \succcurlyeq R_B)$ Strict increase of deductive power

Algebraic Sets Deductive Hierarchy



Algebraic Sets Deductive Hierarchy



Darboux-based

So far we have only seen algebraic sets: p = 0

- For $\wedge_i p_i = 0$, we can rewrite using $\sum_i p_i^2 = 0$
- We can do better [SAS'14]

What about semi-algebraic sets:

- Sets of the form $p \le 0$
- Closed Sets
- Arbitrary Sets: Boolean formulae with polynomial equalities and inequalities

(Non-strict) Barrier Certificate ✓

$$(\forall \mathbf{x} \in \mathbb{R}^n, \mathfrak{D}_\mathbf{f}(p) \leq 0) \longrightarrow p \leq 0$$
 is a positive invariant

Unsound Barrier Certificate X

$$(\forall \mathbf{x} \, s.t. \, \mathbf{p}(\mathbf{x}) = \mathbf{0}, \mathfrak{D}_{\mathbf{f}}(p) \leq 0) \longrightarrow p \leq 0$$
 is a positive invariant

Strict Barrier Certificate ✓

$$(\forall \mathbf{x} \ s.t. \ p(\mathbf{x}) = 0, \mathfrak{D}_{\mathbf{f}}(\mathbf{p}) < \mathbf{0}) \longrightarrow p \leq 0$$
 is a positive invariant

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Conservative Lifting of Barrier Certificates to Boolean Connectives

Differential Invariants (DI)

$$(\mathsf{DI})\frac{D(S)^{\mathbf{f}}_{\dot{\mathbf{x}}}}{S \to [\dot{\mathbf{x}} = \mathbf{f}] \ S},$$

 $D(S)_{\dot{\mathbf{x}}}^{\mathbf{f}}$: substitute each $\dot{\mathbf{x}}_i$ in D(S) with $\mathbf{f}_i(\mathbf{x})$

$$\begin{array}{lll} D(r) & = 0 & \text{for numbers,} \\ D(x) & = \dot{x} & \text{for variables,} \\ D(a+b) & = D(a)+D(b), \\ D(a\cdot b) & = D(a)\cdot b+a\cdot D(b), \\ D(a\leq b) & \equiv D(a)\leq D(b), \text{ accordingly for } \geq,>,<. \text{ (BC)} \\ D(S_1\wedge S_2) & \equiv D(S_1)\wedge D(S_2), \\ D(S_1\vee S_2) & \equiv D(S_1)\wedge D(S_2), \text{ (\wedge here is important for soundness)} \end{array}$$

Liu, Zhan, Zhao (LZZ) Characterization [EMSOFT'2011]

Definitions

$$\begin{aligned} & \operatorname{In}_f(S) \equiv \{\mathbf{x} \in \mathbb{R}^n \mid \exists \ \epsilon > 0. \ \forall \ t \in (0, \epsilon). \ \mathbf{x}(t) \in S\}, \\ & \operatorname{In}_{(-f)}(S) \equiv \{\mathbf{x} \in \mathbb{R}^n \mid \exists \ \epsilon > 0. \ \forall \ t \in (0, \epsilon). \ \mathbf{x}(-t) \in S\}, \end{aligned}$$

Characterization

An arbitrary set $S \subset \mathbb{R}^n$ is a positive invariant for \mathbf{f} if and only if

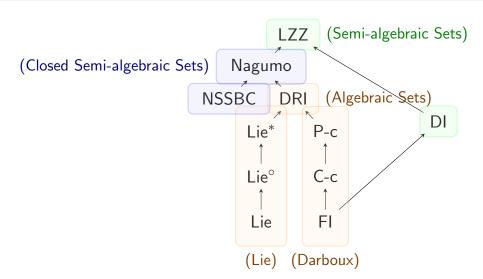
$$S \subseteq \operatorname{In}_f(S)$$
 and $S^c \subseteq \operatorname{In}_{(-f)}(S^c)$

or equivalently

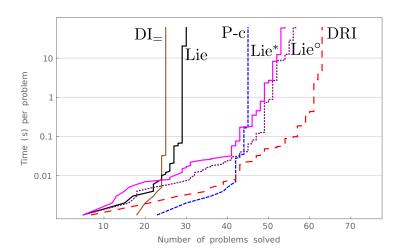
$$\operatorname{In}_{(-f)}(S) \subseteq S \subseteq \operatorname{In}_f(S)$$

• $In_f(S)$ can be computed using high-order Lie derivatives!

Proof Rules for Checking Invariance of Semi-algebraic Sets



Experimental Performance: Algebraic Sets



Template-Based

- Fix a template: generic polynomials with symbolic coefficients
- Use sufficient conditions to derive constraints over the coefficients
- Solve the system
- In practice:
 - templates of the form p=0 (algebraic sets) with Darboux condition $(\mathfrak{D}_{\mathbf{f}}(p) \in \langle p \rangle)$ work for $n \leq 10$ and $d \leq 3$.
 - templates of the form $p \le 0$ relying on Barrier Certificates (SOSTools)

Discrete Abstraction

- Discrete the space with the signs of a given set of polynomials
- Relies on sufficient conditions to remove spurious transitions
- Issue: How to populate and eventually refine the initial set of polynomials

Suppose we have a 2-dimensional ODE $(\dot{x}_1, \dot{x}_2) = (x_1, x_2)$

Suppose we have a 2-dimensional ODE
$$(\dot{x}_1,\dot{x}_2)=(x_1,x_2)$$

1 Start with parametric h of degree 1: $h = \alpha_1 x_1 + \alpha_2 x_2 + \alpha_3$

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$$\mathfrak{D}_{\mathbf{f}}(h) = \alpha_1 x_1 + \alpha_2 x_2 = \beta(\alpha_1 x_1 + \alpha_2 x_2 + \alpha_3)$$

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- Max dim of ker of $M(\beta) \rightsquigarrow$ more freedom for $\alpha = (\alpha_1, \alpha_2, \alpha_3)$
- Increases the chances of finding first integrals
- Dually, minimize the rank of $M(\beta)$

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- Dually, minimize the rank of $M(\beta) \rightsquigarrow NP$ -hard [Buss et al. 1999]

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- Max dim of ker of $M(\beta) \rightsquigarrow$ more freedom for $\alpha = (\alpha_1, \alpha_2, \alpha_3)$
- Increases the chances of finding first integrals
- Dually, minimize the rank of $M(\beta) \rightsquigarrow NP$ -hard [Buss et al. 1999]

$$h = x_2(0)x_1 - x_1(0)x_2$$

Simulation (Time Travel)

A peek in the future (or the past) of a concrete model for a concrete initial (or final) condition. Essentially by "executing" the model step by step: only a **Local** recipe is needed.

Verification (Time Abstraction)

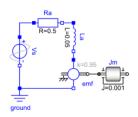
Qualitative analysis of the geometry (shape) of the state space as it captures **Global** properties.

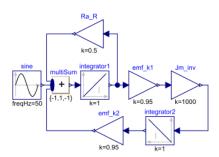
Local description defines the global properties which in turn benefit numerical approximations (eg. Geometrical Integration)

Outline

- 1 Introduction
- 2 Verification
- 3 Simulation
- 4 Challenges

Block Diagram vs. State Flow





Differential-Algebraic Equations (DAE)

 Verification and Faithful Simulation for cyber-physical systems with DAE

Differential-Algebraic Equations (DAE)

 Verification and Faithful Simulation for cyber-physical systems with DAE



- $\mathbf{0} = \mathbf{F}(\mathbf{x}, \dot{\mathbf{x}}, t)$
- $\begin{array}{l}
 \bullet \\
 \mathbf{0} = \mathbf{g}(\mathbf{y}, \mathbf{x})
 \end{array}$
- Compositional design
- Tools: Dymola (Dassault Systèmes), Modelica

if Guard do Differential Equation

- Guard: predicate in the state variables and their time derivatives.
- **Differential Equation**: equation, **implicit** or explicit, in the state variables and their time derivatives

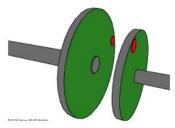
if Guard do Differential Equation

- Guard: predicate in the state variables and their time derivatives.
- Differential Equation: equation, implicit or explicit, in the state variables and their time derivatives.

When a guard holds, its equation is enforced.

Ideal Clutch





if t do
$$J_1\dot{\omega}_1 = \tau_1$$
 (e₁)
if t do $J_2\dot{\omega}_2 = \tau_2$ (e₂)
if γ do $\omega_1 - \omega_2 = 0$ (e₃)
if γ do $\tau_1 + \tau_2 = 0$ (e₄)
if $\neg \gamma$ do $\tau_1 = 0$ (e₅)
if $\neg \gamma$ do $\tau_2 = 0$ (e₆)

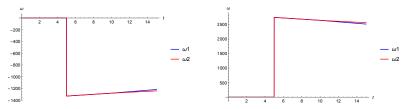
- State Variables: the angular velocities ω_1 and ω_2
- ullet γ is an input signal modelling the pedal's position

Clutch Disengaged, $\gamma = f$ Ordinary Differential Equation

Clutch Engaged, $\gamma = t$ Differential Algebraic Equation

if t do
$$J_1\dot{\omega}_1 = \tau_1$$
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if $\neg \gamma$ do $\tau_2 = 0$ (e₆)

- Dymola crashes with a division by zero
- Mathematica treats resets as initializations (nondeterministic behavior)



The solution may be discontinuous when $\gamma:f\to t$ because of the additional constraint $\omega_1-\omega_2=0$

Problem 1 How to handle overdetermined systems?

- The angular velocities ω_1 and ω_2 are **known**
- ullet γ switches to t (the driver engages the clutch)

$$\omega_1 - \omega_2 = 0$$
 is **enforced**

- The system becomes overdetermined
- The solution is not smooth and even discontinuous

Problem 2 What is the meaning of the derivatives?

Some equations must hold for $\gamma = t$ and $\gamma = f$.

if t do
$$J_1\dot{\omega}_1= au_1$$
 (e₁)
if t do $J_2\dot{\omega}_2= au_2$ (e₂)

- What is the **meaning** of derivatives when $\gamma : f \rightarrow t$?
- How to compute the reset values ?

Solution for Overdetermined Systems

[Benveniste, Caillaud, G., HSCC 2017]

Causality Principle

The additional constraints are

- caused by (consequence of) the current status, and
- enforced at the immediate next instant.

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t: present

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 $\omega_1(t+\delta) - \omega_2(t+\delta) = 0$
 $t+\delta$, $0 < \delta << 1$; immediate future

[Benveniste, Caillaud, G., HSCC 2017]

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 $\delta \in {}^{\star}\mathbb{R}$ is a positive infinitesimal

Nonstandard reals, Hyperreals

Infinite Sequence of Reals

- $\delta = \langle \delta_1, \delta_2, \dots \rangle$
- $\delta_i \in \mathbb{R}$
- Not necessarily convergent
- $\langle 1, \frac{1}{2}, \frac{1}{3}, \dots \rangle$ is a (positive) infinitesimal
- $r = \langle r, r, r, \dots \rangle, r \in \mathbb{R}$
- Functions over the reals can be internalized
- $x(\langle t_1, t_2, \dots \rangle) = \langle x(t_1), x(t_2), \dots \rangle$

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Nonstandard reals, Hyperreals

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Nonstandard Difference Quotient

Let $\delta \in {}^{\star}\mathbb{R}$ be a non zero infinitesimal.

$$\frac{x(t+\delta)-x(t)}{\delta}$$

Proposition

A real function x is differentiable at t if and only if there exists a real number b such that

$$\frac{x(t+\epsilon)-x(t)}{\epsilon}\sim b$$

for any non zero infinitesimal ϵ .

Derivatives as Difference Quotients

$$\dot{x}$$
 is replaced by $\frac{x(t+\delta)-x(t)}{\delta}=\frac{x^{\bullet}-x}{\delta}$

- Shift forward (when needed)
- Formal substitution of time derivatives into difference quotient.

Solving Nonstandard Systems

$$\begin{array}{lll} \text{if t do} & J_1 \frac{\omega_1^\bullet - \omega_1}{\delta} = \tau_1 & (e_1^\delta) \\ \\ \text{if t do} & J_2 \frac{\omega_2^\bullet - \omega_2}{\delta} = \tau_2 & (e_2^\delta) \\ \\ \text{if } \gamma & \text{do} & \omega_1^\bullet - \omega_2^\bullet = 0 & (e_3^\bullet) \\ \\ \text{if } \gamma & \text{do} & \tau_1 + \tau_2 = 0 & (e_4) \\ \\ \text{if } \neg \gamma & \text{do} & \tau_2 = 0 & (e_6) \end{array}$$

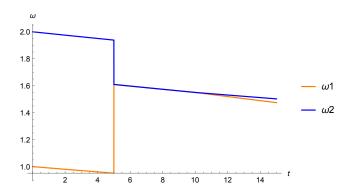
Solving Nonstandard Systems

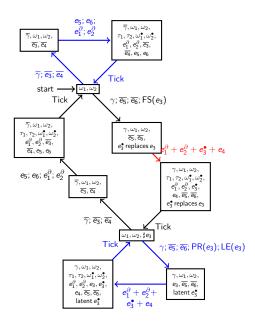
$$\omega_1^{\bullet} = \omega_2^{\bullet} = \frac{J_1 \omega_1 + J_2 \omega_2}{J_1 + J_2}$$

Standardization

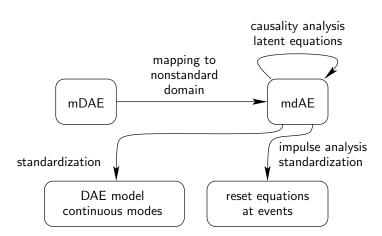
- Automated procedure for a class of systems
- Generalization remains a challenge

Expected Simulation





Unifying Discrete and Continuous Dynamics



Challenges



Combining Static Analysis and Symbolic Computation

Invariants Generation

- Extends previous work (algebraic approach)
- Approximate exact computations to scale

Simulation of multi-mode DAE

- Well-founded operational semantics (compilation, index reduction)
- Preserve composionality in presence of mode changes
- Proper handling of zero-crossing (detection and consistent initialization)
- Cascades of zero-crossings (sliding modes)
- Investigate the use of static analysis and symbolic computation

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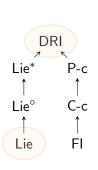
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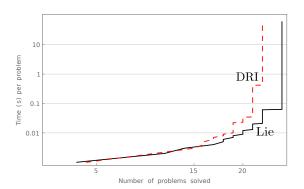
Thanks for your attention!



Smooth invariant manifolds (Lie vs DRI)

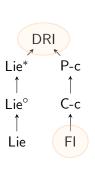
Lie and DRI decide invariance for smooth invariant manifolds.

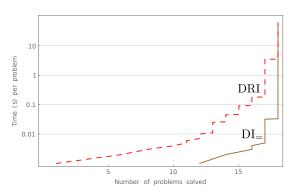




Functional invariants (DI vs DRI)

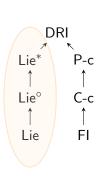
DI= and DRI decide invariance of varieties of conserved quantities.

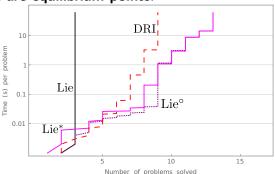




Singularities at Equilibria (Lie, Lie^o & Lie^{*} vs DRI)

Lie°, Lie* and DRI decide invariance for varieties of with singularities that are equilibrium points.





$$(h_1 > \max(h_2, \dots, h_m) \to \mathfrak{D}_{\mathbf{f}}(h_1) < 0)$$

$$\land (h_1 < \max(h_2, \dots, h_m) \to \mathfrak{D}_{\mathbf{f}}(\max(h_2, \dots, h_m)) < 0)$$

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$$(g_1 < \min(g_2, \dots, g_m) \to \mathfrak{D}_{\mathbf{f}}(g_1) < 0)$$

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